application note

powerful adaptec
chip set simplifies
disk controller design

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1.0 INTRODUCTION

To take advantage of the advances in CPU technology, designers need a high performance rigid disk controller that can effectively break the I/O bottleneck. Unlike disk controllers in the past, new implementations have to be flexible and powerful. In fact, today's design must be capable of being an intelligent I/O subsystem, instead of being a drain on the host CPU. In addition, the disk controllers need to be easily adaptable to recent enhancements in the Winchester disk drive technology.

To be flexible, a controller design must be capable of not only handling the slower drive interfaces such as the ST506 and SA1000, but also enhancements like the ST412HP (High Performance) and other high performance interfaces like ESDI and SMD. In addition, the controller must be able to meet not only the original specifications but also various deviations from them. The controller therefore has to deal with variations in track density, track formatting, data transfer rates, sector sizes and many other things.

The Adaptec disk controller chip set provides the major portion of the hardware necessary to build a high performance rigid disk controller. Three custom LSI circuits make up the set: a single chip controller, an encoder-decoder chip, and a buffer controller. The AIC-100 disk controller chip handles high-speed sequencing of the data, leaving slower drive-control operations to an inexpensive support microprocessor. This approach enables the use of the same chip set with all of the industry-standard disk interfaces (ST506, SA1000, ESDI, SMD and ST412HP).
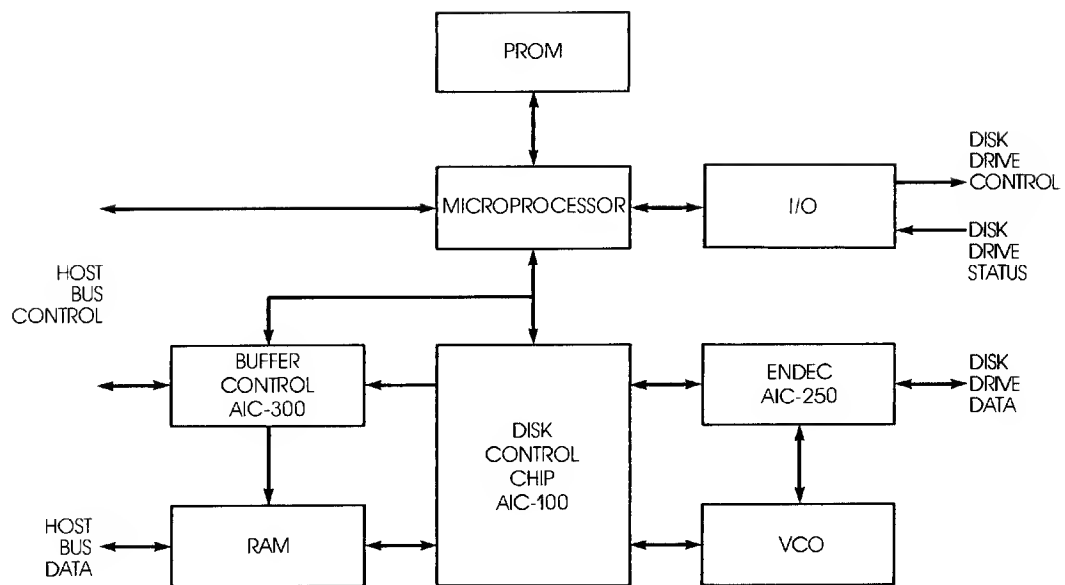
The AIC-250 LSI encoder-decoder circuit along with a data separator forms all that is necessary for NRZ (non return-to-zero) to MFM (modified frequency modulation) conversion and back. The data separator includes the phase locked loop and the voltage controlled oscillator that are necessary to separate the data from the clock during an MFM read data operation. The chip also performs the address mark generation, address mark detection, and write precompensation functions.

In most systems with shared DMA channels, adequate cycle times are not available to transfer data to and from the disk every 800 nsecs per byte (10 Mbits/sec designs) or every 1.6 μsecs per byte (5 Mbits/sec designs). Without some protection built into the board, even a slow Winchester (ST506) could overrun the DMA with some regularity. The AIC-300 dual ported buffer controller can address up to 64K Bytes of RAM buffer in a true dual-ported buffer configuration that guarantees non-interleaved disk formatting and no host overruns. In addition, the chip greatly simplifies the implementation of the popular SCSI (Small Computer System Interface) bus.

A typical controller block diagram is shown in Figure 1.

By partitioning the function between the different chips, the Adaptec disk controller solution addresses many major growth areas in the disk drive market. Some of the extremely important parameters are:

- **Number of Head Selection Lines:** To give way to the demand for increased number of heads in the next generation of drives, this is controlled by a low cost support processor.

- **Seek Commands:** To handle a variety of requirements from step-pulse line for seek control (in ST506 type interfaces) to structured serial output commands (such as direct addressing in ESDI type interfaces), this function is again best handled by the support processor.

```
                          ┌──────────┐
                          │   PROM   │
                          └────┬─────┘
                               ↕                              DISK
                               │                              DRIVE
              ┌────────────────┴─────┐      ┌──────────┐   ►CONTROL
     ◄────────┤   MICROPROCESSOR    ◄├─────►│   I/O    │
              └──────┬───────────────┘      └──────────┤   ►DISK
                     │        ▲                         │    DRIVE
   HOST              │        │                              STATUS
   BUS               │        │
 CONTROL             │        │
         ┌───────────┴──┐  ┌──┴───────────┐   ┌──────────┐   DISK
  ◄──────┤   BUFFER     │  │              │◄─►│  ENDEC   │ ►DRIVE
         │   CONTROL    │  │     DISK     │   │  AIC-250 │◄─ DATA
         │   AIC-300    │  │   CONTROL    │   └──────────┘
         └──────┬───────┘  │    CHIP      │
                │          │   AIC-100    │
   HOST         │          │              │
   BUS ◄─┬──────┴──┐       │              │   ┌──────────┐
   DATA  │   RAM   │◄─────►│              │◄─►│   VCO    │
         └─────────┘       └──────────────┘   └──────────┘
```

Figure 1. Typical Controller Block Diagram

■ **Hard or Soft Sector Formatting:**   The new combination of fixed and removeable drives require hard sector recognition, a function that is handled by the AIC-100 controller chip. A requirement for digital data processing in graphics and sound extends sector lengths to as much as a full track, also supported by the AIC-100.

■ **Placement of Encoding or Decoding Functions:**   Drive manufacturers have taken different approaches to the handling of data. Some drive interfaces require MFM data while others require NRZ data. By separating the NRZ from/to MFM conversions to the AIC-250 ENDEC chip, the AIC-100 disk controller can be used in a variety of controller designs.

■ **Error Detection and Correction:**   All data correction should be transparent to the host, so as not to waste valuable processing time. The AIC-100 controller chip offers the best placement for this function.

■ **Data Buffer Control:**   This function can be easily yet economically handled by using the AIC-300 buffer controller to convert inexpensive RAM into a truly dual ported FIFO. The size of the buffer can be based on the application, up to a maximum of 64K bytes.

This Application Note discusses the use of the Adaptec disk controller chip set in designing a high performance Winchester disk controller.

Section 2 describes the data transfer between the host and the controller. The performance benefits of using the AIC-300 dual port controller will be discussed.

Section 3 describes the use of the AIC-100 disk controller chip in performing the formatting and high speed sequencing functions.

Section 4 discusses the use of a support processor in monitoring the operations of the AIC-100 and the AIC-300. In addition, the implementation of some popular drive interfaces will also be described.
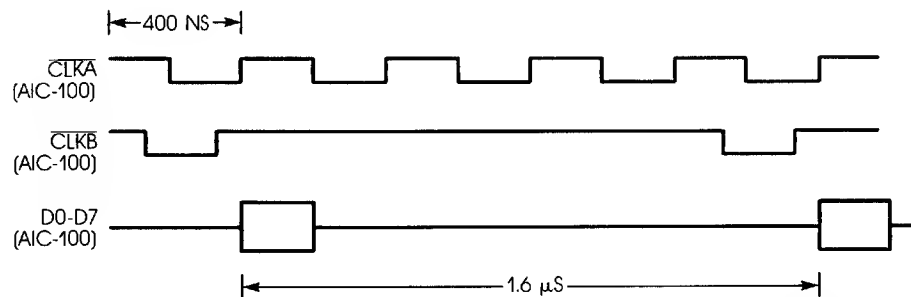
Section 5 deals with the encoding/decoding of the serial data. Here, the use of the AIC-250 ENDEC chip in the NRZ from/to MFM conversion will be discussed.

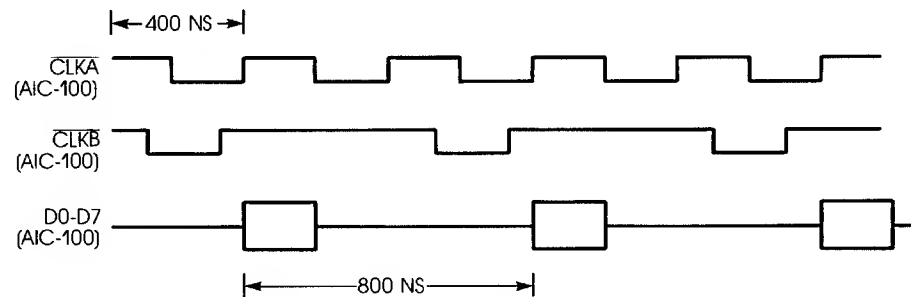Section 6 is a summary, and sums up this Application Note.

Section 7 is the Appendix, and gives some additional detail on the use of the Adaptec chips in the design of a high performance controller.

2.1 PORT A TRANSFER

The effectiveness of a controller design is based on its ability to transfer data to and from the disk as fast as physically possible, and as accurately through Error Detection and Correction (EDAC). Since the drive is continuously spinning during a read or write operation, a byte has to be transferred from/to the AIC-100 disk controller chip every 800 nsec (10 MHz operation) or every 1.6 μsecs (5 MHz operation). The AIC-100 chip either indicates the availability of a byte (during disk read) or requests a byte (during disk write), once every byte time through a signal output on the CLKB line (pin 3). This is shown in Figure 2.

Figure 2a. AIC-300 Port A Data Byte Transfer Timing—5 MHz

Figure 2b. AIC-300 Port A Data Byte Transfer Timing—10 MHz

This data byte is transferred between the AIC-100 and the RAM buffer and is called a Port A transfer. The appropriate address in the RAM is generated by the AIC-300 buffer controller chip, from a set of pointer registers present in the chip. The interface between the AIC-100 disk controller chip and the AIC-300 buffer controller chip necessary to accomplish this transfer is very simple, and is shown in Figure 3.

Figure 3. Port A Transfer Interface

2.2 PORT B TRANSFER

The data transfer between the RAM buffer and the host is initiated and controlled by the AIC-300 buffer controller chip, and is referred to as the Port B transfer. The AIC-300 generates the necessary request signal (Port B Req), RAM buffer address and control signals, and the host bus latch control signals. This is shown in Figure 4. The buffer controller uses the Port B Request (pin 39) to initiate the transfer which is completed after a Port B Acknowledge is received (sampled at pin 38).

Figure 4. Port B Transfer Interface

The relationship between the different control signals during a Port B transfer are shown in Figure 5.

Figure 5a. AIC-300 Port B Data Transfer—Host to Buffer

Figure 5b. AIC-300 Port B Data Transfer—Buffer to Host

2.3 DATA TRANSFER OVERVIEW

In a controller design, the AIC-300 provides all the necessary signals to interface to the host and the disk controller chip. The buffer controller also generates the necessary control signals to access the buffer, alternating between Port A and Port B transfers. While the Port A transfer is synchronous in nature (at the data transfer frequency), the Port B transfer is asynchronous and is based on a Req/Ack handshake with the host. Figure 6 shows the timing overview.

Figure 6a. Dual Port Buffer Timing Overview—5 MHz

Figure 6b. Dual Port Buffer Timing Overview—10 MHz

Thus during any data transfer operation, even though the access to the buffer is interleaved between Port A and Port B, the AIC-100 disk controller chip should stay one sector ahead of the host. This is necessary to keep error detection and correction transparent to the host. The AIC-300 has an internal stop pointer which is used to prevent overruns. An overview of the buffer operation is shown in Figure 7.

Figure 7. Buffer Operation Overview

Thus during a read operation, while data is being transferred into the buffer from the drive (Port A transfer), at the same time, the previous sector can be transferred to the host (Port B transfer). The stop pointer is set to the end of the sector being transferred to the host, and prevents an overrun. At the conclusion of a successful sector read, the stop pointer can be updated to point to the end of this sector. Now this sector can also be transferred to the host, while yet another sector is read in.

2.4   RAM BUFFER INTERFACE

As mentioned earlier, the buffer controller is capable of managing up to a 64K byte buffer. The chip provides the address, memory select and read/write lines necessary to perform this function. There are two distinct modes of RAM interface: equal to or less than 1K bytes (up to 10 address lines), and greater than 1K bytes (up to 16 address lines).

In the 10 bit or less addressing mode, two special lines supply A8 and A9 address lines. An example is shown in Figure 8.

Figure 8. 10-Bit Addressing Application Example

In the 16 bit addressing mode (greater than 10), the higher order address lines (A8-A15) and the lower order address lines (A0-A7) are multiplexed coming out of the chip, on pins A0 to A7. In this mode, the two special lines supply the strobes SDP and SHP used to demultiplex these signals during a disk transfer (port A) and a host transfer (port B) respectively. An example is shown in Figure 9.

Figure 9. 16-Bit Addressing Application Example

Internal to the AIC-300 are two sets of pointer registers used to generate the RAM addresses. These are the Read Address Pointer (RAP) and the Write Address Pointer (WAP). The actual register used depends on which port is selected and the direction of the data transfer, and is controlled by the value of the ROP/WOP (Read Operation/Write Operation) bit in the DMA control register (Reg 53, bit 4). In addition to this, another pair of registers called the Stop Pointer (SP) is used to control data transfers between the host and the buffer. The different possibilities are shown in Table 1.

Table 1. Buffer Address Generation

ROP/WOP	*Source for Address*	*Direction of Transfer*
Read Disk		
1	WAP	Disk to Buffer
1	RAP	Buffer to Host (RAP < SP)
Write Disk		
0	RAP	Buffer to Disk
0	WAP	Host to Buffer (WAP < SP)

3.1   DRIVE INTERFACE

The AIC-100 Winchester disk controller chip handles all the data formatting and sequencing necessary to interface with the drive.

During a write operation, parallel data is transferred into the chip from the RAM buffer. This data is passed through a 32 bit shift register before it is output as serial data in the NRZ format. During this time, a 4 byte ECC is also computed. After the entire sector is written to the drive, the ECC bytes are also output in a serial fashion.

During a read operation, NRZ data is read into the controller chip and byte parallel data is transferred to the RAM buffer. Error checking is also performed at this time, using the four ECC bytes stored at the end of the sector. If there is an error, the support processor usually initiates retries to determine if the error is correctable. If the error is determined to be correctable, the error syndrome information found in the AIC-100 chip is used to correct the data in the buffer. This process is transparent to the host.

The serial data is synchronized through a read reference clock that is input to the AIC-100 chip. The disk controller chip has the necessary logic to look for index mark and sector mark (on hard sectored drives), in addition to generating the Read Gate and Write Gate signals. The disk controller also controls the writing of the address mark during a write operation, and looks at when the address mark is detected during a read operation. Figure 10 shows the necessary drive interface lines.

Figure 10.   AIC-100 Drive Interface Signals

3.2 REGISTER OVERVIEW

Internally, the AIC-100 disk controller is made up of 17 registers, used to control the high speed sequencing functions. The slower functions such as seek, head selection, drive selection, and drive status monitoring are handled by an external support microprocessor and some I/O port logic. This is described in the following section.

In addition to the 17 internal registers, the AIC-100 chip also decodes accesses to 5 external register addresses which can be used to simplify controller design. The registers can be grouped by function as follows:

- Stored Value Registers
- Command Ste-up Registers
- Command Register, Register 78
- Status/Execution Register, Register 79
- External Registers

Table 2 shows the registers in each group and their function.

The operation of the AIC-100 chip revolves around two of these registers: the command register (Reg 78), and the status register (Reg 79). These registers have to be set up and monitored by the support microprocessor for all data sequencing operations.

Table 2.

Register	Title	R=Read/W=Write	Function
		STORED VALUE REGISTERS	
72	ECC Bits 0-16 "OR"ed and 17-23	R	Syndrome bits.
73	ECC Bits 24-31	R	Error bits.
D0	Gap Length	W	Bit control.
E0	Cylinder Byte	R/W	ID field byte.
E1	Head Byte	R/W	ID field byte.
E2	Sector Byte	R/W	ID field byte.
E3	Flag Byte	R/W	
		COMMAND SET-UP REGISTERS	
71	ECC Control	W	Correction control.
74	ECC Polynomial	W	Low order bits.
77	ECC Polynomial	W	High order bits.
7A	Operation Modifier	R/W	Operation control.
7E	Special I/O	R	Input and data transfer bits.
7F	POP Stack	R	LIFO stack read.
7F	Clock Control	W	CLKA control.
A4	Search Bit	W	Enables search.
C4	Sector Length	W	Sector data field length.
		COMMAND REGISTER	
78	Operation Command	W	Controls the sequencing of the controller chip.
		STATUS/EXECUTION REGISTERS	
79	Chip Status	R	Used to monitor chip status.
79	Start Execution	W	Used to start chip execution.
		EXTERNAL REGISTERS	
50	Host Data Transfer	R/W	Allows processor to R/W data directly from the host.
51	Host Data Transfer	R/W	Exactly like Register 50.
70	Buffer Data Transfer	R/W	Allows processor to R/W data directly from the buffer.
6E	External Line Control	R/W	Causes the pins designated RD6E and WR6E to be activated by the processor.
6F	External Line Control	R/W	Same as above, but for pins RD6F and WR6F.

4.0 SUPPORT PROCESSOR

4.1 PROCESSOR INTERFACE

The AIC-100 chip and the AIC-300 chip are setup and monitored by a support processor. The interface to the support microprocessor is through a multiplexed address/data bus as is found in the Intel 8085 family of microprocessors. This can however be easily adapted to other processors, such as the Z80, through the use of minimal external logic. Some examples of this are shown in Appendix 7.5. For the rest of this controller design application note, an Intel 8085 is used. The basic interface between the 8085 and the Adaptec chips is shown in Figure 11.

Figure 11. Support Processor Interface

The support processor is used to maintain "loose" synchronization with what is happening in real-time on the disk through the OP command and Sequencer Status registers in the AIC-100 (Reg 78 and Reg 79). Based on the operation, the processor also sets up the registers in the AIC-300 to control the data transfer to and from the buffer.

The sequence of the actual steps that have to be followed with the AIC-100 and the AIC-300 for a READ, WRITE, and FORMAT disk operation is shown in the form of flow charts in Appendix 7.1. The actual source code is shown in Appendix 7.2.

4.2 DISK CONTROL AND STATUS MONITORING

An important function of the support processor is the management of the low speed drive interface function. By designating this operation to the support processor, the Adaptec chip set can be easily adapted to a variety of drive interfaces and drive types.

In the case of a ST506/412 interface, the support processor has to perform the following functions:

- Recalibrate to Track 000.
- Monitoring Drive Ready and Write Fault.
- Selecting the appropriate drive and head.
- Seeking to the desired cylinder.

This can be easily accomplished through the use of I/O ports and firmware. The AIC-100 chip internally decodes accesses to addresses 6E and 6F. A read or write to these two addresses cause the appropriate pin (pin 4, 5, 6, or 7) on the AIC-100 chip to be asserted. Thus by placing external I/O ports at addresses 6E and 6F, and using this feature of the AIC-100 chip, the designer can eliminate external decode logic. An example of this implementation for a ST506/412 or a ST412HP interface is shown in Figure 12.

Since the support processor is used for drive control, a variation in hardware and a change in the firmware enables the Adaptec chip set to be used for other drive interfaces. An example of the implementation for an SMD interface is shown in Figure 13.

Figure 12. ST506/412 and ST412HP Interface

Figure 13. SMD Drive Interface

The diagram shows the following connections:

- **8085** connected via **MICRO BUS** to **244**
- **244** has inputs: FAULT, SEEK ERROR, ON CYLINDER, UNIT READY, WRITE PROTECTED, BUSY; output labeled E
- **373** (upper) outputs: UNIT SEL TAG, UNIT SEL 0, UNIT SEL 1, TAG 1, TAG 2, TAG 3, BIT 8, BIT 9
- **373** (lower) outputs: BIT 0, BIT 1, BIT 2, BIT 3, BIT 4, BIT 5, BIT 6, BIT 7
- **AIC-100** with inputs/outputs: AD (0–7), RD, WR, ALE, $\overline{R6E}$ (5), $\overline{W6E}$ (4), $\overline{W6F}$ (6), INDEX (10), SECTOR (11)
- SECTOR and INDEX lines at bottom

4.3 HOST BUS ACCESS (REGISTERS 50 AND 51)

In the design of intelligent controllers, a popular host interface used is the SCSI (Small Computer System Interface), previously known as the SASI. Adaptec chip set based designs ease the implementation of the SCSI, since a majority of the functionality needed for this is inherent in the overall architecture and, in particular, in the AIC-300 chip.

While the AIC-300 buffer controller chip can handle data transfers between the host bus and the buffer, it may sometimes be necessary for the support processor to directly access the host data bus. This is the case in an SCSI implementation during the transfer of the Command Data Block (CDB) and message bytes. This can be accomplished by the support processor through an access to Register 50 or 51.

An access to either register is decoded by both the AIC-100 and the AIC-300. The AIC-100 chip internally bridges the support processor data bus and the RAM buffer data bus, thus offering a data path to host latch and receivers. The AIC-300 chip asserts the control signals necessary to access the host bus. During a read, the BIE line from the AIC-300 is asserted. During a write, the LO line will be first asserted, followed by the BOE line. This allows the data to be latched first before being enabled onto the host data bus. During a read or a write to Register 50 or 51, the data is passed through the AIC-100. An example is shown in Figure 14.

Figure 14. Host Bus Access

4.4 RAM BUFFER ACCESS (REGISTER 70)

During the process of transferring data between the disk and the host, the support processor has to sometimes have access to the sector data. This is especially necessary during a read operation if a correctable ECC error is encountered. Using the error

syndrome information found in the AIC-100 chip, the support processor first calculates an error mask and a displacement. The algorithm necessary for this is shown in Appendix 7.3.

After that, the actual data stored in the buffer has to be read, modified, and written back. This is done through an access to Register 70.

Here again the AIC-100 decodes an access to Register 70 and bridges the processor data bus and the buffer data bus. This allows the transfer of data between the support processor and the RAM.

The AIC-300 decodes an access to Register 70 and asserts memory select and read/write to the RAM. The address selected is the contents of the WAP registers if the ROP/WOP bit (Reg 53, bit 4) is set (Read Disk) and the contents of the RAP registers if the ROP/WOP bit is reset (Write Disk). This path is shown in Figure 15.

Figure 15. RAM Buffer Access

5.0 | **SERIAL DATA ENCODING/DECODING**

The AIC-100 controller chip converts parallel data to NRZ serial data and back. In drive interfaces, such as SMD and ESDI, serial data is written to, and read from, the drive interface in NRZ format. However, in drive interfaces like the ST506/412, SA1000, and the ST412HP, the data is transferred in MFM format. In this case, the main difference is the transfer rate, shown in Table 3.

Table 3. MFM Transfer Rates

Interface	Transfer Rate
SA1000	4.34 Mbits/s
ST506/412	5.0 Mbits/s
ST412HP	10.0 Mbits/s

The AIC-250 Encoder/Decoder chip (ENDEC) contains a major portion of the logic necessary to convert data in the NRZ format to/from MFM format. In addition, the chip provides the logic necessary for address mark detection and writing. The chip also has the necessary circuitry for write precompensation which is necessary to compensate for certain bit patterns.

During a read disk operation, an external VCO/PLL is required to separate the data from the clock before it can be converted to NRZ data. A typical VCO/PLL circuit is shown in Appendix 7.4.

An example of the ENDEC interface circuit is shown in Figure 16.

The AIC-250 simplifies ST506/412 and ST412HP interface Winchester disk controller designs due to the fact that it needs very few external passive components, unlike other implementations which require costly delay lines. A typical configuration of the ENDEC for a 5 MHz data rate is shown in Figure 17.

Figure 16. NRZ To/From MFM Conversion

Figure 17. External Component Requirement at 5 MHz Data

6.0 | SUMMARY

As with all growing technologies, 5¼" Winchester drives will undoubtedly move beyond the capacity of today's units. Thus a designer who is designing a controller board today must consider the needs of tomorrow. To that end, by cleverly partitioning features among various chips in the Adaptec family, and by putting the burden of slow speed disk control and status monitoring on the support processor, future design issues can also be addressed. All of this of course being accomplished without sacrificing performance or increasing cost.

| 7.0 | **APPENDICES** |

7.1 FLOWCHARTS

In order to read or write data from/to the peripheral, the support processor has to loosely control the operation of the AIC-100 Winchester disk controller chip and the AIC-300 dual ported buffer controller chip. The flowcharts necessary to set up and monitor these two chips are as follows:

- AIC-100
  - □ Soft Sector Format
  - □ Soft Sector Read/Write
    (Single and Multi Block)
- AIC-300
  - □ Read Disk (Single and Multi Block)
  - □ Write Disk (Single and Multi Block)

```
                              LOAD CHSF          E0–E3  ┐
                                                        │
                         LOAD GAP 1 AND          D0     │
                         GAP 3 LENGTH                    │
                                                        │
                    LOAD BLOCK COUNTER           C4     │  SET UP
                    80H=256; 00H=128                     │
                                                        │
                 LOAD OP CMD WITH 10H=           78     │
                 WRITE GAP 1 AND ID                      │
                                                        │
                    SET B5 IN 7A FOR             7A     ┘
                    6CH PATTERN

                    LOAD START  15H              79     ┐
                                                        │
                      READ STATUS                79     │  WAIT FOR
                                                        │  INDEX
                  NO      INDEX                          │
                         R79 B5                          ┘

                          YES  ◄──── (1)

                    LOAD OP CMD                  78     ┐
                    11H=WRITE DATA                       │
                                                        │
                      READ STATUS                79     │  WRITE ID
                                                        │
                  NO     DATA                            │
                        TRANSFER                         │
                        R79 B6                           ┘
                                      VARIABLE
                          YES  ◄────  SECTOR
                                      LENGTH
                                      ROUTINE
```

AIC-100 Soft Sector Format

LOAD OP CMD 13H=FORMAT TO END OF TRACK — 78 — NO — MORE BLOCKS — YES — LOAD CHSF — E0–E3

READ STATUS — 79 — LOAD OP CMD 12H=WRITE ID — 78

WRITE 4E UNTIL INDEX — NO — STOPPED R79 B4 — READ STATUS — 79

YES — END

NO — BRANCH ACTIVE R79 B5 — YES — (1) — WRITE DATA

AIC-100 Soft Sector Read/Write (Single and Multi Block)

Single Block

```
CLEAR ALL
POINTERS
REG 59=0
```
⎤
⎥ COMMON
⎥ READ
⎥ STEPS
⎦
```
SET ROP
REG 53 BIT 4
```

```
START DATA XFER
FROM PERIPHERAL
TO BUFFER
```

IS XFER DONE — NO

YES

```
SET SP=(WAP-1)
```

```
START DMA
SET RL
REG 53 BIT 3
```

IS DMA DONE REG 53 BIT 5 — NO

YES

END

Single Block

Multi Block

DO COMMON READ STEPS

IS BUFFER AVAILABLE FOR NEXT BLOCK — NO → XFER TO HOST TO MAKE ROOM FOR NEXT BLOCK

YES

```
START DATA
XFER FROM
DEVICE
```

IS XFER DONE — NO

YES

```
SET SP=(WAP-1)
```

MORE BLOCKS TO XFER — YES

NO

IS DMA DONE REG 53 BIT 5 — NO

YES

```
CLEAR
READ LATCH
```

END

Multi Block

AIC-300 Read Disk

Single Block

Multi Block

AIC-300 Write Disk

7.2 SAMPLE ROUTINES

In order to control the operation of the disk controller, the support processor has to execute firmware that loosely synchronizes the Adaptec chip set. This section gives sample routines necessary to accomplish some basic functions. These can be modified to meet the specific design requirement. The following flowchart gives an overview of the disk control operation.

Disk Controller Operation Overview

The following is a breakdown of the main functional sections in the disk controller operation.

- **Power On Reset**:   Initialize all known parameters for both host and drive interfaces. Clear or set appropriate flags to indicate unknown parameters (i.e., dirve not formatted). Some initial command may be necessary from the host in order to fully specify all drive or host parameters.

- **Idle Loop**:   Constantly check the host interface for incoming commands. Do any other housekeeping or checking necessary to remain active and alert to the host or drive interface. Seek complete on one of the drives may be checked here.

- **Command Decode**:   Check for incoming commands from the host for correct formats, lengths, and zero check unused fields. Save the host I.D. and the logical unit number of the active unit. Possibly separate command types here and do the seek, if necessary. Jump to the routine to execute the command.

- **The Commands**:   Start and stop command execution, ensure data transfers to or from the host, check I.D. and data fields for validity, and maintain proper positioning on the drive. Please refer to the sample routine listings that follow this text for guidelines.

- **Final Status**:   Report to the host upon completion of each command. Report proper completion or error conditions. Clean up host and drive interfaces, internal flags, etc., and return to the idle loop.

Some typical error conditions are as follows:

Drive Errors	Controller Errors	System Error
NO INDEX SIGNAL	I.D. ERROR	INVALID COMMAND
NO SEEK COMPLETE	UNCORRECTABLE DATA ERROR	ILLEGAL BLOCK ADDRESS
WRITE FAULT	I.D. ADDRESS MARK NOT FOUND	VOLUME OVERFLOW
DRIVE NOT READY	DATA ADDRESS MARK NOT FOUND	BAD ARGUMENT
TRACK 00 NOT FOUND	RECORD NOT FOUND	INVALID LUN
	SEEK ERROR	
	CORRECTABLE DATA ERROR	
	INTERLEAVE ERROR	
	UNFORMATTED DRIVE	
	SELF TEST ERROR	
	DEFECTIVE TRACK	

Following are some sample routines for AIC-300 and AIC-100 chips. These routines show how to program normal disk functions. Error routines are left to the user to define. The drive interface defined here may also be different for some users. The chip registers are defined by their locations for easy understanding without having to refer to the equate list.

```
****************************************************************
*                    SYSTEM EQUATES                           *
****************************************************************

                 BUFFER CHIP

R53     EQU     4053H           ;DMA control register
R54     EQU     4054H           ;BUFFER capacity
R59     EQU     4059H           ;RESET
R5A     EQU     405AH           ;RAP register
R5C     EQU     405CH           ;WAP register
R5E     EQU     405EH           ;STOP register

                 SERDES CHIP

R70     EQU     4070H           ;BUFFER I/O DATA
R71     EQU     4071H           ;ECC control register
R72     EQU     4072H           ;ECC low byte
R73     EQU     4073H           ;ECC high byte
R74     EQU     4074H           ;ECC poly byte FORWARD
R77     EQU     4077H           ;ECC poly byte RECIPROCAL
R78     EQU     4078H           ;BRANCH CONTROL register
R79     EQU     4079H           ;START register
R7A     EQU     407AH           ;OPERATION CONTROL register
R7F     EQU     407FH           ;PUSH/POP I.D. stack/CLOCK DIVIDE
RE0     EQU     40E0H           ;CYLINDER area
RE1     EQU     40E1H           ;HEAD area
RE2     EQU     40E2H           ;SECTOR area
RE3     EQU     40E3H           ;FLAG area
RC4     EQU     40C4H           ;DATA field length
RA4     EQU     40A4H           ;SEARCH OPERATION cell
RD0     EQU     40D0H           ;GAP length cell
ECCOFF  EQU     7               ;ECC HARDWARE OFFSET
****************************************************************
*         RESET BUFFER AND SERDES CHIPS                       *
****************************************************************
RESET:  DI                      ;disable interrupts
        MVI     A,20H           ;reset serdes
        STA     R71             ;send to serdes
        MVI     A,00H           ;RESET OFF
        STA     R71             ;send to serdes
        STA     R71             ;and again
        MVI     A,1             ;reset...
        STA     R59             ;BUFFER CHIP
        XRA     A               ;RESET OFF
        STA     R59
        MVI     A,42H
        OUT     080H            ;INITIALIZE PORTS
        JMP     IDLE            ;GO TO IDLE LOOP
****************************************************************
*          FORMAT ENTIRE DRIVE ROUTINE                        *
****************************************************************
FORMAT: MVI     B,00            ;clear flag
        MVI     A,40H
        STA     R53             ;ENABLE INITIATOR
        CALL    GETDRV          ;SELECT THE DRIVE
        CALL    RECAL           ;recal drive
CLRLA:  XRA     A               ;zero out
        STA     RE0             ;cylinder...
        STA     RE1             ;head...
        STA     RE2             ;sector...
        STA     RE3             ;and flag
        MVI     A,0FFH
        STA     RC4             ;SET 256 BYTE SECTORS
        MVI     A,20H           ;supress xfer on
        STA     R7A             ;send to op control reg
NEWTRK: MVI     C,32            ;32 sectors per track
        MVI     A,0EH           ;gap 1
        STA     RD0             ;store in serdes
        MVI     A,10H           ;set up value...
        STA     R78             ;for branch register
        LXI     H,R79           ;point H/L at start register
        MVI     M,15H           ;write start value
BRACT0: MOV     A,M             ;get status
        ANI     20H             ;test branch active bit
        JZ      BRACT0          ;loop til branch active OR INDEX
        MVI     A,11H           ;WRITE I.D. branch value
        STA     R78             ;save in branch register
        ANI     40H             ;test xfer active
```

```
            JZ      XFACT0          ;loop til it is
            MVI     A,0CH           ;gap 3
            STA     RD0             ;save in serdes
            DCR     C               ;decrement sector counter
            JZ      LSTSEC          ;if last sector, jump out
            MVI     A,12H           ;otherwise, use 12...
            STA     R78             ;as branch
            CALL    INCLA           ;GET NEXT SECTOR/STORE NEW C,H,S,F
            JMP     BRACT0          ;and loop back to branch active check
LSTSEC: MVI A,13H                   ;stop value
            STA     R78             ;send to branch register
LSTWT0. MOV A,M                     ;get SERDES status
            ANI     10H             ;test the stopped bit
            JZ      LSTWT0          ;loop till stopped AT INDEX
        USERS MUST DEFINE THEIR OWN ROUTINES INCLA AND BUMP.
        INCLA IS INTERNAL HOUSEKEEPING ROUTINE TO KEEP TRACK
            OF CYLINDER AND HEAD UNTIL ENTIRE DISK IS FORMATTED.
            IT ALSO STORES NEW C,H,S,F IN RE0 - RE3.
        BUMP WILL SELECT A NEW HEAD AND SEEK, WHEN NECESSARY.


            CALL    INCLA
            CALL    BUMP
            JNC     NEWTRK          ;if not done, loop back
FMDONE: XRA A                       ;if done, zero out...
            STA     R7A             ;supress bit
            JMP     IDLE            ;and jump to idle loop
*****************************************************************
*           READ ONE OR MULTIPLE BLOCKS ROUTINE         *
*****************************************************************
READ:   XRA A                       ;zero out...
            MOV     D,A             ;CLEAR FOR MSB STOP POINTER
            STA     R59             ;reset ALL POINTERS on BUFFER CHIP
            MVI     A,0D0H          ;set READ OPERATION
            STA     R53             ;send to BUFFER CHIP
            CALL    GETDRV          ;select drive
            LDA     NBLKS           ;GET BLOCK COUNT
            MOV     C,A             ;save in C
            LDA     CYL             ;GET CYLINDER
            STA     RE0             ;INTO SERDES
            STA     DESTRK          ;save as desired track
            LDA     HEAD            ;GET HEAD
            STA     RE1             ;INTO SERDES
            STA     DESTHD          ;save as desired head
            CALL    SEEK            ;do the seek
            JC      BADCMD          ;on error, halt
            XRA     A               ;zero out..
            STA     RE3             ;FLAG
            LDA     SECT            ;get sector
            STA     RE2             ;INTO SERDES
RDRENT: MVI A,08H                   ;stuff 08
            STA     R78             ;into branch register
            LXI     H,R79           ;point H/L at start register
            MVI     E,03            ;3 revs to find i.d.
            LDA     R7A             ;dummy read to clear index passed
RDSTRT: MVI M,05H                   ;START SERDES
RDDLY:  MOV A,M                     ;get status
            MOV     B,A             ;save a copy in B
            ANI     20H             ;check branch active
            JNZ     RDDAM           ;branch active = GOOD I.D.-READ DATA
            MOV     A,B             ;recover status
            ANI     10H             ;check stopped bit
            JZ      RDDLY           ;if not stopped, loop back up
            MOV     A,B             ;if stopped, get new copy
            ANI     04H             ;look for ecc error
            JNZ     RDIDER          ;if ecc blown, jump to error routine
            LDA     R7A             ;if not ecc stop, get op ctl
            ANI     01              ;test index passed
            JZ      RDSTRT          ;if not passed, loop back
            DCR     E               ;if passed, drop rev counter
            JNZ     RDSTRT          ;TRY AGAIN
            JMP     SEEKER          ;if zero, ERROR
RDDAM:  MVI B,0                     ;CLEAR COUNTER
            ANA     A
            JM      RDXFLP          ;TEST FOR DATA TRANSFER ACTIVE
            LDA     R79
            ANA     A
            JM      RDXFLP
            DCR     B               ;INSURE NO HANG IF NO DATA A.M.
            JNZ     AMNACT          ;WAIT FOR DATA TRANSFER
            JMP     DAMERR          ;NO DATA A.M. FOUND
```

```
RDXFLP: MVI     A,5             ;USE 05
        STA     R78             ;as branch
RDNXT:  DCR     C               ;DECREMENT COUNTER
        JZ      RDBR2           ;if zero, get out
        CALL    INCLA           ;GET NEXT SECTOR
RDBR2:  MOV     A,M             ;get status
        MOV     B,A             ;save a copy in B
        ANI     30H             ;test branch active and stopped bits
        JZ      RDBR2           ;loop till branch active or stopped
        MVI     A,08H           ;get the next branch address into A
        STA     R78             ;and save in branch register
        MOV     A,B             ;now recover status
        ANI     05H             ;test for ecc or compare error
        JNZ     RDDAER          ;jump out on error
        LDA     R53             ;DATA IS GOOD
        ANI     08H             ;CHECK READ LATCH
        JZ      RD4             ;NOT ON - FIRST SECTOR
        LDA     R53
        ANI     20H             ;CHECK DMA DONE
        JNZ     RD2             ;DMA IS DONE - SEND NEXT SECTOR
        MVI     1FH
        STA     R79             ;DMA NOT DONE - STOP SERDES
RD1:    LDA     R53
        ANI     20H
        JZ      RD1             ;WAIT FOR DMA DONE
        INR     A               ;BUMP IT
        CPI     4               ;INSURE NOT OVER 3
        JNZ     RD3
        XRA     A               ;BACK TO 0
RD3:    MOV     D,A             ;SAVE IT
        STA     R5F             ;PUT IT IN STOP HIGH
RD4:    MVI     A,0FFH
        STA     R5E             ;SET STOP LOW = FF
        MVI     A,0D8H
        STA     R53             ;SET READ LATCH AND START DMA
        MOV     A,C             ;CHECK BLOCK COUNT
        ANA     A
        JZ      STOPRD          ;DONE
        .MP     RDRENT          ;GO READ NEXT SECTOR
STOPRD: MVI     A,14H           ;stop location...
        STA     R78             ;into branch register
STOPBR: MOV     A,M             ;GET STATUS
        ANI     10H             ;look at stopped bit
        JZ      STOPBR          ;loop till it is
        MOV     A,M             ;get status again
        ANI     04H             ;test ecc error bit
        JNZ     RDDAER          ;if set, jump to ecc retry routine
STOP1:  LDA     R53             ;GET STATUS
        ANI     20H
        JZ      STOP1           ;WAIT FOR DMA DONE
        XRA     A
        STA     R53             ;TURN OFF READ LATCH
        CALL    KILSER          ;STOP SERDES
        JMP     IDLE            ;otherwise, back to idle loop
************************************************************
*           WRITE ONE OR MULTIPLE BLOCKS ROUTINE         *
************************************************************
WRITE:  XRA     A               ;zero out...
        MOV     D,A             ;CLEAR FOR MSB STOP POINTER
        STA     R59             ;BUFFER CHIP RESET
        CALL    GETDRV          ;select drive
        LDA     NBLKS           ;GET BLOCK COUNT
        MOV     C,A             ;save in C
        LDA     CYL             ;GET CYLINDER
        STA     RE0             ;INTO SERDES
        STA     DESTRK          ;save as desired track
        LDA     HEAD            ;GET HEAD
        STA     RE1             ;INTO SERDES
        STA     DESTHD          ;save as desired head
        CALL    SEEK            ;do the seek
        JC      BADCMD          ;on error, halt
        XRA     A               ;zero out...
        STA     R59             ;buffer pointers
        XRA     A               ;zero out..
        STA     RE3             ;FLAG
        LDA     SECT            ;get sector
        STA     RE2             ;INTO SERDES
WRRENT: MVI     A,09H           ;stuff 09 (start address)...
        STA     R78             ;into branch register
```

```
        LXI     H,R79           ;point H/L at start register
        MVI     E,03            ;3 revs to find i.d.
        LDA     R7A             ;dummy read to clear index passed
        MOV     A,D             ;GET MSB STOP POINTER VALUE
        INR     A               ;BUMP IT
        CPI     4               ;INSURE NOT OVER 3
        JNZ     WD1
        XRA     A               ;BACK TO 0
WD1:    MOV     D,A             ;SAVE VALUE
        STA     R5F             ;SET STOP POINTER HIGH
        MVI     A,0FFH
        STA     R5E             ;STOP POINTER LOW = FF
        MVI     A,0C4H
        STA     R53             ;SET WRITE LATCH - START DMA
WD2:    LDA     R53             ;GET STATUS
        ANI     20H
        JZ      WD2             ;WAIT FOR DMA DONE
WRSTRT: MVI     M,05H           ;DATA IS IN BUFFER - START SERDES

WRDLY:  MOV     A,M             ;get status
        MOV     B,A             ;save a copy in B
        ANI     20H             ;check branch active
        JNZ     WRDAM           ;if branch active, jump
        MOV     A,B             ;recover status
        ANI     10H             ;check stopped bit
        JNZ     WRRENT          ;if stopped, re-enter
        LDA     R7A             ;if not stopped, get op ctl
        ANI     01              ;test index passed
        JZ      WRDLY           ;if not passed, loop back
        DCR     E               ;if passed, drop rev counter
        JNZ     WRDLY           ;TRY AGAIN
        JMP     SEEKER          ;if zero, find out why
WRDAM:  MOV     A,M             ;get status
        MOV     B,A             ;save a copy in B
        ANI     10H             ;stopped ??
        JNZ     WRRENT          ;if stopped, re-enter to re-start
        MOV     A,B             ;recover status
        ANI     40H             ;transfer active ??
        JZ      WRDAM           ;if no, loop back

WRXFLP: MVI     A,05H           ;use 05...
        STA     R78             ;as branch
        XRA     A               ;zero out...
        STA     WRFLAG          ;write flag

WRNXT:  DCR     C               ;decrement counter
        JZ      STOPWR          ;if zero, get out
        CALL    INCLA           ;GET NEXT SECTOR
        MVI     D,09H           ;next branch (09) into D to tighten code

WRBR2:  MOV     A,M             ;get status
        MOV     B,A             ;save a copy in B
        ANI     30H             ;test branch active and stoped bits
        JZ      WRBR2           ;loop till branch active or stopped
        MOV     A,D             ;get the next branch address into A
        STA     R78             ;and save in branch register
        MOV     A,B             ;now recover status
        ANI     10H             ;check stopped bit
        JNZ     WRRENT          ;re-enter if stopped
        JMP     WRDAM           ;otherwise, loop back up
STOPWR: MVI     A,14H           ;STOP LOCATION
        STA     R78             ;into branch register
WRSTBR: MOV     A,M             ;get status
        ANI     10H             ;look at stopped bit
        JZ      WRSTBR          ;loop till it is
        XRA     A
        STA     R53             ;TURN OFF WRITE LATCH
        CALL    KILSER          ;STOP SERDES
        JMP     IDLE            ;otherwise, back to idle loop
```

```
*****************************************************
*       RECAL ROUTINE DOES A ONE-STEP-AT-        *
*       A-TIME RECAL WHILE LOOKING FOR THE       *
*       TRACK Ø SIGNAL.                          *
* NOTE: USERS INTERFACE MAY BE DIFFERENT THAN    *
*       SHOWN BELOW FOR PIN ASSIGNMENTS FOR      *
*       STATUS, SEEK PULSES AND SEEK DIRECTION.  *
*****************************************************
RECAL:  PUSH    PSW             ;temp save A/PSW
        MVI     A,01H           ;select drive 1, direction=out
        OUT     82H             ;send to port
RDYCK1: IN      83H             ;get status
        ANI     08H             ;test ready
        JZ      RDYCK1          ;loop till ready
TRØØCK: IN      83H             ;get status
        ANI     02H             ;test trk Ø
        JNZ     RETØ            ;if track Ø, get out
        MVI     A,41H           ;step pulse
        OUT     82H             ;send it
        MVI     A,01H           ;un-step
        OUT     82H             ;send it
SKCPLP: IN      83H             ;get drive status
        ANI     01H             ;look at seek complete
        JZ      SKCPLP          ;loop till seek complete
        JMP     TRØØCK          ;then check for track ØØ
RETØ:   XRA     A               ;zero out
        STA     CURTRK          ;current track
        POP     PSW             ;recover A/PSW
        RET                     ;return
*********************************************************
*       SEEK TO TRACK AND HEAD IN DESTRK AND DESTHD,   *
*       WHICH CONTAIN CYLINDER AND HEAD, RESPECTIVELY. *
*       USER TO SET THESE UP FROM COMMAND DATA.        *
*********************************************************
SEEK:   PUSH    PSW             ;temp save A/PSW
        PUSH    B               ;temp save B/C
        LDA     CURTRK          ;get current track
        MOV     B,A             ;save in B
        LDA     DESTRK          ;get destination track
        MOV     C,A             ;save in C
        CMP     B               ;compare them
        JZ      DOHEAD          ;if they are the same, jump around
        JC      GOOUT           ;if desired is less than current, jump
GOIN:   MOV     A,C             ;get destination track
        SUB     B               ;subtract smaller current track
        MOV     C,A             ;put result in C
        IN      82H             ;get control port
        ORI     80H             ;set dir in
        JMP     DOSEEK          ;and jump to common code
GOOUT:  MOV     A,B             ;get current track
        SUB     C               ;subtract smaller destination track
        MOV     C,A             ;put result in C
        IN      82H             ;get control port
        ANI     7FH             ;set dir out
DOSEEK: OUT     82H             ;set up direction
SEEKLP: ORI     40H             ;set step pulse
        OUT     82H             ;send it
        ANI     ØBFH            ;un-step
        OUT     82H             ;send it
WTSKC1: IN      83H             ;get drive status
        ANI     01H             ;look at seek complete
        JZ      WTSKC1          ;loop till seek complete
        DCR     C               ;drop seek counter
        JNZ     SEEKLP          ;loop till C=Ø
        LDA     DESTRK          ;get derired track
        STA     CURTRK          ;save as current track

        NOTE:   USERS INTERFACE MAY BE DIFFERENT THAN SHOWN BELOW
                FOR DRIVE AND HEAD SELECTION.

DOHEAD: LDA     DESTHD          ;get desired head
        RAL                     ;shift left...
        RAL                     ;to put head number...
        RAL                     ;in correct position...
        RAL                     ;to send to I/O port
        ANI     30H             ;zap any garbage
        MOV     B,A             ;save in B
        IN      82H             ;get current control port
        ANI     ØCFH            ;clear head select
```

```
                ORA    B                ;OR in head
                OUT    82H              ;output result
                POP    PSW              ;recover A/PSW
                RET                     ;and return
        ****************************************************
        *              SELECT DRIVE ROUTINE               *
        ****************************************************
        GETDRV: PUSH   PSW              ;temp save A/PSW
                IN     82H              ;get current port value
                ANI    ØFCH             ;zap select bits
                ORI    Ø1               ;select drive
                OUT    82H              ;send it out
                POP    PSW              ;recover A/PSW
                RET                     ;and return


        *********************************************************************
        *       STOPS THE SERDES AND THEN RETURNS                          *
        *********************************************************************
        KILSER: PUSH   PSW              ;temp save A/PSW
                MVI    A,1FH            ;then load the stop location..
                STA    R79              ;to the start register to halt serdes
        HLCKLP: LDA    R79              ;get serdes status
                ANI    1ØH              ;check the stopped bit
                JZ     HLCKLP           ;loop till halted
                POP    PSW              ;if stopped, recover A/PSW
                RET                     ;and then return
        *********************************************************************
        *              ROUTINE READS THE SYNDROME INTO ECC1                *
        *********************************************************************
        RDSYND: MVI    A,Ø4H            ;DISABLE FEEDBACK
                STA    R71              ;TO SERDES
                LXI    H,R73            ;point H/L at serdes ecc register
                LXI    D,ECC1           ;point D/E at ram to save syndrome
                MOV    A,M              ;get msb of syndrome
                STAX   D                ;save in ram
                INX    D                ;bump ram pointer
                MVI    C,Ø3             ;three more bytes to do
        BYTLOP: MVI    B,Ø8             ;8 bits per byte
                MVI    A,Ø6H            ;inhibit feedback and shift value
        SHFT:   STA    R71              ;force shift of syndrome
                DCR    B                ;subtract 1 from bit counter
                JNZ    SHFT             ;loop 8 times (1 byte)
                MOV    A,M              ;load the next byte of the syndrome
                STAX   D                ;save it in ram
                INX    D                ;and bump ram pointer
                DCR    C                ;drop byte count
                JNZ    BYTLOP           ;loop 3 times
        ****************************************************
        *         ATTEMPT TO CORRECT ECC ERROR             *
        ****************************************************
        FIXIT:  MVI    A,Ø8H            ;CLEAR ECC COMMAND
                STA    R71              ;TO SERDES
                MVI    A,Ø4H            ;RESET OFF
                STA    R71
                LXI    H,ECC1+3         ;H/L = FIRST SYNDROME BYTE
                MVI    C,Ø4             ;TO ENTER 4 BYTES
        NXTBYT: MVI    B,Ø8             ;8 BITS PER BYTE
                MOV    D,M              ;GET BYTE TO LOAD
        SHFBYT: MOV    A,D              ;BYTE TO LOAD IN A
                RAR                     ;LSB BIT TO CARRY
                MOV    D,A              ;SAVE SHIFTED BYTE
                RAL                     ;LSB BIT TO BIT Ø
                ANI    Ø1               ;ONLY ONE BIT
                ORI    Ø6H              ;ADD SHIFT COMMAND
                STA    R71              ;STUFF BIT VIA ECC CONTROL REG
                DCR    B                ;DECREMENT BIT COUNT
                JNZ    SHFBYT           ;LOOP 8 TIMES
                DCX    H                ;TO NEXT HIGHER SYNDROME BYTE
                DCR    C                ;DECREMENT BYTE COUNT
                JNZ    NXTBYT           ;LOOP 4 TIMES
        STFRCP: LXI    H,REVERS         ;POINT TO REVERSE POLYNOMIAL
                CALL   LDPOLY           ;INTO POLYNOMIAL REGISTERS
        FIGLEN: LHLD   ABLKSZ           ;GET CURRENT BLOCK SIZE
                PUSH   H                ;MOVE POINTER
                POP    B                ;INTO B/C
                MVI    A,Ø8             ;NUMBER OF BITS PER BYTE
                CALL   MULT             ;NUMBER OF BITS PER BLOCK
                LXI    H,3CH            ;4 ECC + 2 GAP BYTES TIMES 8
                DAD    B                ;H/L = TOTAL BITS PER BLOCK
```

```
                SHLD    TEMP            ;SAVE TOTAL BIT COUNT
                XCHG                    ;PUT IN D/E
                MVI     A,00            ;DISABLE ECC SHIFTS
                STA     R71             ;TO ECC CONTROL REG
                LXI     H,R72           ;H/L = ERROR TEST REG
SHFTST: MVI     A,02H           ;TO SHIFT ECC
                STA     R71             ;TO CONTROL REG
                DCX     D               ;DECREMENT BIT COUNT
                MOV     A,M             ;GET ERROR INDICATOR
                ANA     A               ;TEST FOR ZERO
                JZ      GOTDIS          ;0 = ERROR FOUND
                MOV     A,E             ;LSB BIT COUNTER
                ORA     D               ;CHECK FOR ZERO
                JNZ     SHFTST          ;NON ZERO = SHIFT AGAIN
                JMP     ECCERR          ;DONE HERE = HARD ERROR
GOTDIS: CALL    COMPD           ;D/E = MINUS COUNT
                LHLD    TEMP            ;H/L = TOTAL BIT COUNT
                DAD     D               ;TOTAL BITS MINUS COUNT DOWN
                LXI     D,-ECCOFF       ;ECC HARDWARE OFFSET
                DAD     D               ;SUBTRACT HARDWARE OFFSET
                LXI     D,-32           ;FIRST 32 = ECC OR GAP
                DAD     D               ;SUBTRACT 32
                JNC     FIXED           ;ERROR IN ECC = DATA O.K.
                SHLD    TEMP1           ;SAVE COUNT
                MOV     B,03            ;TO DIVIDE BYTE ADDRESS BY 8
RR3LOP: MOV     A,H             ;HIGH BYTE OF BIT COUNT
                ANA     A               ;CLEAR CARRY
                RAR                     ;LSB BIT TO CARRY
                MOV     H,A             ;SAVE ROTATED VALUE
                MOV     A,L             ;GET LOW BYTE
                RAR                     ;CARRY TO BIT 7
                MOV     L,A             ;SAVE SHIFTED LOW BYTE
                DCR     B               ;DECREMENT SHIFT COUNTER
                JNZ     RR3LOP          ;LOOP 3 TIMES
                MOV     A,H             ;HIGH BYTE OF SHIFTED ADDRESS
                ANI     1FH             ;CLEAR 3 HIGH BITS
                MOV     H,A             ;SAVE IT IN H
                XCHG                    ;SHIFTED ADDRESS TO D/E
                CALL    COMPD           ;D/E = MINUS ADDRESS
                LHLD    R5C             ;UPPER WRITE LIMIT POINTER (WAP)
                DCX     H               ;LAST DATA BYTE
                DAD     D               ;SUBTRACT POINTER FROM TOP OF DATA
CORECT: LDA     R73             ;GET ERROR MASK DATA
                MOV     D,A             ;D = ERROR MASK
                MVI     E,0             ;CLEAR E
                LDA     TEMP1           ;BIT COUNT = ERROR OFFSET IN 2 BYTES
                ANI     07              ;ONLY 3 BITS
                INR     A               ;PLUS ONE = SHIFT COUNT
                MOV     B,A             ;B = SHIFT COUNT
RRXLP:  DCR     B               ;DECREMENT SHIFT COUNT
                JZ      MASKOK          ;0 = MASK SHIFTED TO PROPER BITS
                MOV     A,D             ;GET HIGH BYTE OF MASK
                ANA     A               ;CLEAR CARRY
                RAR                     ;BIT ZERO TO CARRY
                MOV     D,A             ;RESTORE SHIFTED COPY
                MOV     A,E             ;GET LOW BYTE
                RAR                     ;CARRY TO BIT 7
                MOV     E,A             ;RESTORE LOW BYTE
                MOV     A,D             ;RESTORE HIGH BYTE
                ANI     07FH            ;CLEAR BIT 7
                MOV     D,A             ;SAVE HIGH BYTE
                JMP     RRXLP           ;LOOP AGAIN
MASKOK: MOV     A,E             ;GET LOW BYTE OF MASK
                CALL    MIRROR          ;SWAP END FOR END
                MOV     E,A             ;SAVE IN E
                MOV     A,D             ;GET HIGH BYTE
                CALL    MIRROR          ;SWAP END FOR END
                MOV     D,A             ;SAVE IN D
                LDA     R53             ;CURRENT DMA CONTROL BYTE
                ANI     0F7H            ;KILL READ LATCH
                STA     R53
                PUSH    H               ;SAVE ERROR POINTER
                LHLD    R5C             ;GET CURRENT WAP
                SHLD    TEMP            ;SAVE IN MEMORY
                POP     H               ;RESTORE ERROR POINTER
                SHLD    R5C             ;PUT IN WAP
                LDA     R70             ;GET DATA FROM BUFFER
                XRA     E               ;CORRECT LOW BYTE
                STA     R70             ;CORRECTED BYTE TO BUFFER
                INX     H               ;POINTER TO HIGH BYTE
```

```
          SHLD    R5C                 ;PUT IN WAP
          LDA     R70                 ;GET DATA FROM BUFFER
          XRA     D                   ;CORRECT HIGH BYTE
          STA     R70                 ;CORRECTED BYTE TO BUFFER
          LHLD    TEMP                ;RECOVER SAVED WAP
          SHLD    R5C                 ;RESTORE WAP
FIXED:    LXI     F,FORWRD            ;H/L = FORWARD POLYNOMIAL
          CALL    LDPOLY              ;RELOAD FORWARD POLYNOMIAL
          MVI     A,00                ;STANDARD ECC COMMAND
          STA     R71                 ;TO ECC CONTROL REG
          RET


**************************************************************
*         LOAD ECC POLYNOMIAL                               *
**************************************************************
LDPOLY:   PUSH    PSW                 ;H/L = POLY CONSTANT
          PUSH    B                   ;SAVE B/C
          PUSH    D                   ;SAVE D/E
          MVI     B,04                ;TO LOAD 4 BYTES
          LXI     D,R74               ;D/E = FIRST POLY REG
          CALL    MOVBYT              ;LOAD POLYNOMIAL
          POP     D
          POP     B
          POP     PSW
          RET
**************************************************************
*         COMPLEMENT THE D REGISTER                         *
**************************************************************
COMPD:    PUSH    PSW
          MOV     A,D
          CMA
          MOV     D,A
          MOV     A,E
          INX     D                   ;TWO'S COMPLEMENT
          POP     PSW
          RET
**************************************************************
*         MULTIPLY A BY B/C                                 *
**************************************************************
MULT:     PUSH    H
          ANA     A                   ;CLEAR CARRY
SHFTA:    RAR                         ;SHIFT MULTIPLIER
          LNC     SHFTB               ;GO IF NO BIT
          DAD     B                   ;ACCUMULATE FINAL VALUE
          RC                          ;RETURN IF OVERFLOW
SHFTB:    PUSH    PSW                 ;SAVE MULTIPLIER
          MOV     A,C                 ;SHIFT MULTIPLICAND
          RAL
          MOV     C,A
          MOV     A,B
          RAL
          MOV     B,A
          POP     PSW                 ;RESTORE MULTIPLIER
          ANA     A                   ;CHECK IF DONE
          JNZ     SHFTA               ;CONTINUE IF NOT
          MOV     B,H                 ;RESULTS TO B/C
          MOV     C,L
          POP     H
          RET
**************************************************************
*         MIRROR SWAPS ACCUMULATOR END FOR END              *
**************************************************************
MIRROR:   PUSH    H                   ;SAVE H/L
          PUSH    D                   ;SAVE D/E
          MVI     H,8                 ;BIT COUNT = 8
          MOV     D,A                 ;START VALUE INTO D
MIRLOP:   MOV     A,D                 ;GET INITIAL DATA
          RAR                         ;BIT 0 TO CARRY
          MOV     D,A                 ;SAVE SHIFTED VALUE
          MOV     A,E                 ;GET RESULTING VALUE
          RAL                         ;LSB BIT INTO LOW END SHIFTING HIGHER
          MOV     E,A                 ;SAVE RESULTS IN E
          DCR     H                   ;DECREMENT BIT COUNTER
          JNZ     MIRLOP              ;CONTINUE UNTIL DONE
          MOV     A,E                 ;PUT RESULTS IN A
          POP     D
          POP     H
          RET
```

```
FORWRD: DB      00,00,00,00     ;FORWARD ECC POLYNOMIAL
REVERS: DB      00,00,00,00     ;REVERSE ECC POLYNOMIAL

*************************************************************
*                    RAM ASSIGNMENTS                       *
*************************************************************

CURTRK: DB      00H             ;current track

DESTRK: DB      00H             ;destination track

DESTHD: DB      00H             ;destination head

ERRTYP: DB      00H             ;error type

RDFLAG: DB      00H             ;read flag

WRFLAG: DB      00H             ;write flag

CYL:    DB      00H             ;INPUT CYLINDER

HEAD:   DB      00H             ;INPUT HEAD

SECT:   DB      00H             ;INPUT SECTOR

NBLKS:  DB      00H             ;INPUT NUMBER OF BLOCKS

TEMP:   DB      00H             ;TEMPORARY STOGAGE

TEMP1:  DB      00H             ;TEMPORARY STORAGE

ECC1:   DS      4               ;ECC location

ABLKSZ: DS      2               ;BLOCK SIZE

        END
```

7.3 PROGRAMMING 8-BIT ECC CORRECTION

After each read date operation, a read error may have occurred. This may be determined by reading Register 79. If bit 2 is set, an error did occur and the following procedure is employed to determine if the error is correctable. Note that the majority of read errors are soft (i.e., caused by noise) and that the correction algorithm is time consuming. It is recommended that the record be re-read before attempting correction.

The general flow of the algorithm for 8-bit correction is as follows:

1. Off-load the 32-bit syndrome into local RAM.

2. Shift the syndrome back into the ECC register in reverse order, swapping the syndrome end for end.

3. Change the ECC polynomial from forward to reciprocal.

4. Shift the ECC until all bits except the high order (24-31) bits are zero (correctable) or the number of shifts are greater than the number of bits in the record (uncorrectable).

5. If correctable, the number of shifts represent the displacement of the error from the end of the record (the last bit of the ECC). The error pattern is located in bits 24-31 of the ECC register. This pattern is exclusive ORed with the appropriate bits in memory to correct the error.
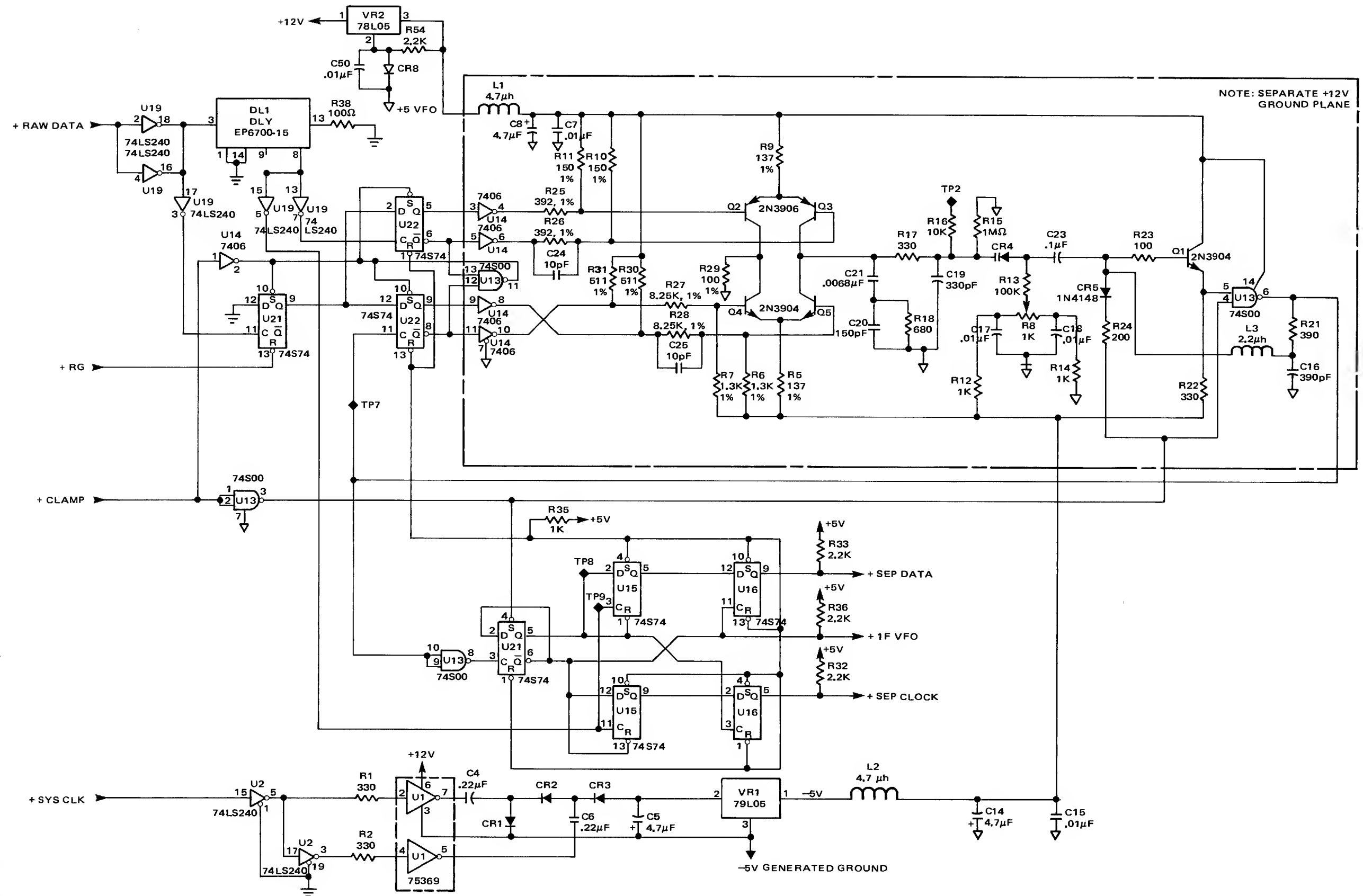
Detailed Programming Steps

1. After a read error is detected, disable feedback by setting R71 = 04_H.

2. Store contents of R73 in RAM (x).

3. Shift ECC 8 times by setting R71 = 06_H eight times.

4. Store contents of R73 in RAM $(x + 1)$.

5. Shift ECC 8 times by setting R71 = 06_H eight times.

6. Store contents of R73 in RAM $(x + 2)$.

7. Shift ECC 8 times by setting R71 = 06_H eight times.

8. Store contents of R73 in RAM $(x + 3)$.

9. Clear ECC and disable feedback by setting R71 to 08 and then 04.

10. Right rotate location RAM $(x + 3)$ and test if carry is set (i.e., test bit 0):
    If set, then load R71 = 07_H.
    If not set, then load R71 = 06_H.
    Repeat operation 7 more times to load entire byte.

11. Repeat step 10 for RAM locations $x + 2$, $x + 1$, and x until all 32 bits of the syndrome are loaded into the ECC in reverse order.

12. Load R74 = 00_H and R77 = 01_H to enable the reciprocal polynomial and disable the forward polynomial.

13. Compute record length in bits:
    # of bits per data field = ECC + Data + AM and SYNC
    For a 256 byte record length in bits = 4 * 8 + 256 * 8 + 2 * 8 = 2096

14. Enable feedback by setting R71 = 00_H.

15. Shift ECC once by setting R71 = 02_H and increment a software counter.

16. Test to see if the software counter is greater than the record length:
     If yes, the error is uncorrectable. Re-enable the forward polynomial and end
     operation.

17. Test to see if R72 = 00$_H$:
     If yes, go to step 18.
     If no, go to step 15.

18. Subtract hardware offset of 7 from the shift count. If a correctable error is located
     within the ECC or the SYNC and AM bytes (the shift count < = 32), the data field
     is good and no further action is required. Subtract 32 from the shift count.

19. The bit displacement (shift count) must now be converted to a byte offset by right
     shifting the count 3 times. The value of the shift count equals the bit displacement
     from end of the record.

20. R73 is the mirror image of the error pattern. Form the error mask data (2 bytes) by
     R73 concatinating with a zero byte.

21. Get the shift count (E#) for error mask data by extracting the lower 3 bits from the
     shift count obtained in step 18.

22. Right shift the error mask data with MSB (bit 15) set to zero. Repeat E#−1 times
     more.

23. Mirror the error mask data byte by byte.

7.4   TYPICAL VCO/PLL SCHEMATIC (DATA SEPARATOR CIRCUIT)

In this section, a typical data separator circuit is shown. This circuit is used in the data
encode/decode portion of the disk controller design. During a read operation, the AIC-
250 chip receives data in the MFM format. In this format, both clock and data informa-
tion are encoded together. The data separator is used to obtain separate data and separate
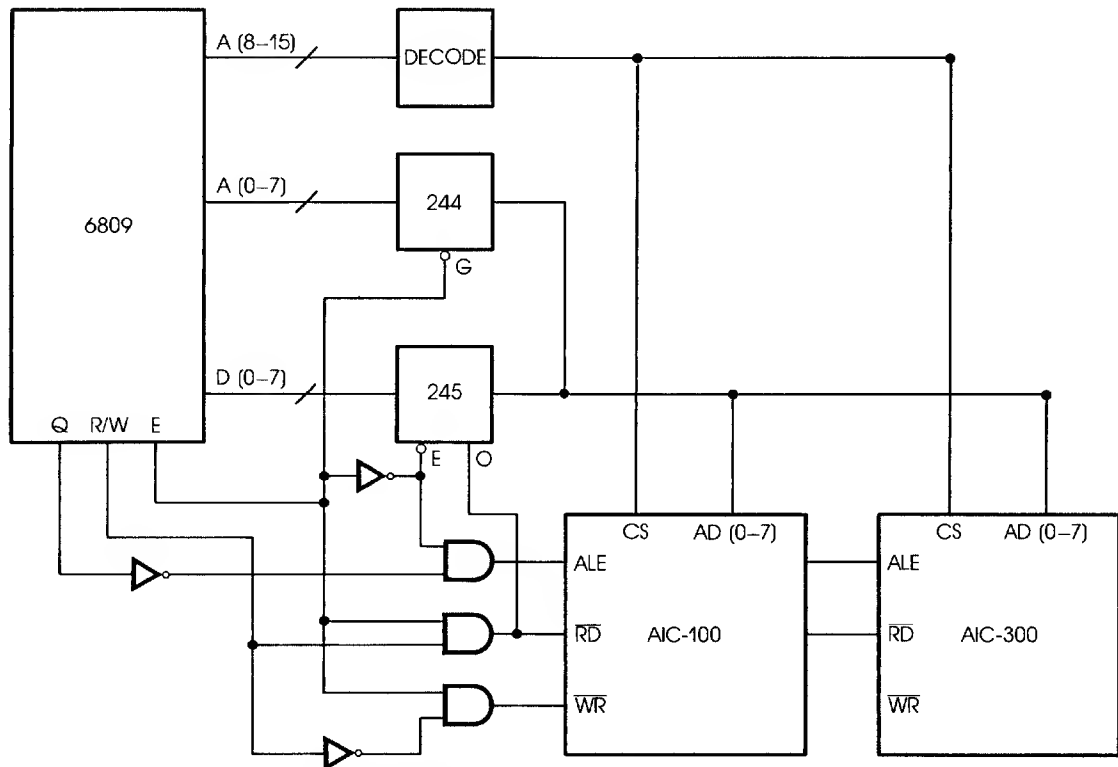clock, from the MFM encoded raw data.

This schematic is commonly used by Adaptec and other disk controller manufacturers.
The user has the choice of using this or any other design.

Typical VCO/PLL Schematic
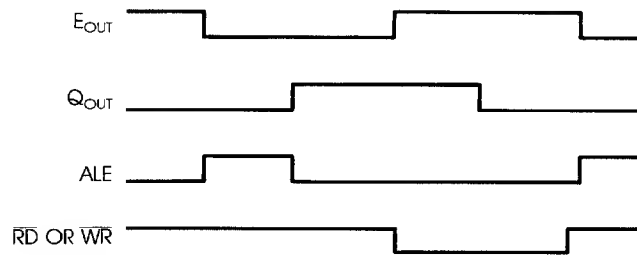
7.5 OTHER MICROPROCESSOR INTERFACES

The Adaptec chip set is intended for use in a system with a multiplexed address/data bus, such as is found in the Intel 8085 family of microprocessors. The NSC 800 can also be very easily used since, in the Z80 emulation mode, the external interface is through an Intel 8085 type multiplexed address/data bus.

Other microprocessors, however, can be easily adapted for use in the disk controller design as the support processor. The basic design involves multiplexing the address and data busses, through external latches, and generating an Address Latch Enable signal (ALE). In addition to this, a separate read and write control signal is required. The following schematic shows the use of a Motorola MC6809 as the support processor in the design of the disk controller.

MC6809 Based Support Processor Interface

In the case of the 6809, as can be seen from the figure, the E_{OUT} and Q_{OUT} signals are used to generate the ALE signal. The generation of RD and WR is fairly simple also. The timing relationships for these signals are shown below.

MC6809 Interface Timing Relationship

In the case of a Z80 microprocessor, the interface is fairly similar. The lower order address and data lines have to be multiplexed in a similar fashion. However, separate RD and WR signals are already available, coming out of the processor chip. The ALE signal can be derived from the IORQ signal which is generated by the Z80 during input or output cycles (IN or OUT class of instructions).